MAMA/215, NST3AS/215, MAAS/215

# MAT3 MATHEMATICAL TRIPOS Part III

Monday 9 June 2025  $-1{:}30~\mathrm{pm}$  to  $4{:}30~\mathrm{pm}$ 

# **PAPER 215**

# MIXING TIMES OF MARKOV CHAINS

### Before you begin please read these instructions carefully

Candidates have THREE HOURS to complete the written examination.

Attempt no more than **THREE** questions. There are **FIVE** questions in total. The questions carry equal weight.

### STATIONERY REQUIREMENTS

Cover sheet Treasury tag Script paper Rough paper

### SPECIAL REQUIREMENTS None

You may not start to read the questions printed on the subsequent pages until instructed to do so by the Invigilator. 1

(a) Consider a Markov chain X on a finite state space S with invariant distribution  $\pi$ . Under which conditions is a randomised stopping time  $\tau$  corresponding to this chain also a strong stationary time?

(b) We consider the following way in which k people shuffle a deck of n cards. One shuffle consists of each person taking a card which is currently on the top of the deck until n - k cards are left in the deck, and then, one by one, in random order, each of the k people places their card in the uniformly random location in the deck (to return their card the first person chooses from n - k + 1 locations, the second person chooses from n - k + 2, and so on until the final person chooses from n locations). Let S be the set of all permutations of  $\{1, \ldots, n\}$ , X the Markov chain on S corresponding to this method of shuffling and  $\pi$  its invariant distribution.

(i) Check that  $\pi$  is the uniform distribution on S.

(ii) Let  $\tau$  be a random time, which is the first time such that at time  $\tau - 1$ , there are at most k - 1 cards above the card which was initially the bottom card of the deck. Explain why  $\tau$  is a strong stationary time.

(iii) For  $k \ll n$ , show that for all  $\varepsilon \in (0,1)$  the mixing time of X, satisfies  $t_{\min}(\varepsilon) \leq \frac{n \log n}{k} + Cn$  for large enough constant C depending on  $\varepsilon$ .

[You can use without proof any results from lectures provided that they are clearly stated. You can use without proof that the mean of a geometric random variable with a parameter q is  $\frac{1}{q}$  while the variance is at most  $\frac{1}{q^2}$  and that  $\sum_{i=1}^{n} \frac{1}{i} = \log n + O(1)$ .

Hint: In part (b)(iii) you may wish to consider some independent Bernoulli trials with success probabilities which differ for different trials and which are all at least p and couple suitably the time it takes to see the first success with a geometric random variable with parameter p.]

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(a) Let X be a Markov chain on a finite state space S with invariant distribution  $\pi$ . For every  $A \subset S$  set  $\Phi(A) = \frac{\sum_{x \in A, y \in A^c} \pi(x) P(x, y)}{\pi(A)}$ .

(i) Define the bottleneck ratio  $\Phi_*$  and the isoperimetric profile  $\Phi_*(r)$  for  $r \leq \frac{1}{2}$ .

(ii) Let  $A = A_1 \cup A_2 \cup \ldots \cup A_k$  for some disjoint sets  $A_i \subset S$  such that for all  $i \neq j$  with  $i, j \leq k$  for all  $x \in A_i, y \in A_j$  we have that P(x, y) = 0. Show that  $\Phi(A) \ge \inf_{i \leq k} \Phi(A_i)$ .

(b) Let  $V_1 = ([0, n^2] \times [0, n^2]) \cap \mathbb{Z}^2$  be the set of points with integer coordinates in a square box of side lengths  $n^2$  and let  $E_1$  be the lattice edges connecting  $(x_1, x_2)$  and  $(y_1, y_2)$  from  $V_1$  if and only if  $|x_1 - y_1| + |x_2 - y_2| = 1$ . Let  $(V_2, E_2)$  be another copy of the pair  $(V_1, E_1)$ . Let  $E_{1,2}$  be a set of edges connecting for all  $k, \ell \in \{0, 1, \ldots, n\}$  the vertex with coordinates  $(kn, \ell n)$ , from  $V_1$ , with the vertex which has exactly the same coordinates in  $V_2$  (there are  $(n + 1)^2$  edges in  $E_{1,2}$ ). Let  $G = (V_1 \cup V_2, E_1 \cup E_2 \cup E_{1,2})$ . Show that the mixing time of lazy, simple random walk on G satisfies  $t_{\text{mix}} \leq n^4$ .

[You are allowed to use any results from lectures without proof. For part (b), it might be useful to recall that for a lazy chain  $t_{\text{mix}} \lesssim \int_{4\pi_{\text{min}}}^{4/M} \frac{du}{u\Phi_*^2(u)} + \frac{1}{\gamma}\log(M)$  where  $\gamma$  is the spectral gap.]

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(a) Suppose that a Markov chain with transition matrix P takes values in a finite vertex set V of a graph G with length function  $\ell$ . Let  $\rho$  be the corresponding path metric,  $\alpha \in \mathbb{R}$  and let

$$\rho_K(\mu, \nu) = \inf \{ \mathbb{E}[\rho(X, Y)] : (X, Y) \text{ is a coupling of } \mu \text{ and } \nu \}$$

be the transportation metric corresponding to  $\rho$ . Suppose that for all edges (x, y) there exists a coupling  $(X_1, Y_1)$  of  $P(x, \cdot)$  and  $P(y, \cdot)$  such that

$$\mathbb{E}_{x,y}[\rho(X_1, Y_1)] \leqslant e^{-\alpha}\rho(x, y).$$

Show that for any  $x, y \in V$  we have  $\rho_K(P(x, \cdot), P(y, \cdot)) \leq e^{-\alpha}\rho(x, y)$  and use this to show that for any probability measures  $\mu$  and  $\nu$  we have  $\rho_K(\mu P, \nu P) \leq e^{-\alpha}\rho_K(\mu, \nu)$ .

Let diam  $(V) = \max_{x,y} \rho(x, y)$  and deduce that

$$t_{\min}(\varepsilon) \leqslant \frac{1}{\alpha}(\log(\operatorname{diam}(V)) + \log(1/\varepsilon)).$$

[You can assume that a transportation metric corresponding to a path metric is a well-defined metric and that it upper bounds the total variation distance.]

(b) Let X be a Markov chain on state space of all subsets of  $\{1, \ldots, n\}$  of size at most k, i.e.  $V = \{A \subset \{1, 2, \ldots, n\} : |A| \leq k\}$ , where k < n, which evolves as follows. If the current state is a set  $A \in V$  we first pick m to be a uniform number from  $\{0, 1\}$  and move the chain depending on the value of m. If m = 0 the chain stays in place. If m = 1 we pick a uniform number i in  $\{1, \ldots, n\}$ . If  $i \in A$  the chain moves to  $A \setminus \{i\}$ , while if  $i \notin A$ , it moves to  $A \cup \{i\}$  provided that |A| < k, and otherwise, it stays in place. For  $A, B \in V$  let  $\rho(A, B) = \frac{|A \setminus B| + |B \setminus A| + ||A| - |B||}{2}$ . Show that  $\rho$  is a path metric corresponding to some length function  $\ell$ . Use it to show that the mixing time of X satisfies  $t_{\min}(\varepsilon) \leq n \log(\frac{k}{\varepsilon})$ .

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In this question, you can use any results from lectures without proof.

(a) For a finite, reversible, Markov chain X on the state space V define the quantity  $\operatorname{hit}_{\alpha}(\varepsilon)$  for  $\alpha, \varepsilon \in (0, 1)$ . Define a  $\operatorname{hit}_{\alpha}$  cutoff. Does mixing time cutoff imply a  $\operatorname{hit}_{\frac{1}{2}}$  cutoff?

(b) Let d > 0 be an integer constant which does not depend on n and let  $G = G_n = (V, E)$  be a graph for which |V| = n is an even number and where each vertex has degree d. Let X be a simple random walk on G, let its mixing time be  $t_{\text{mix}}$  and let its relaxation time  $t_{\text{rel}}$  satisfy  $t_{\text{rel}} \ll t_{\text{mix}}$ . Let  $\theta = \frac{1}{t_{\text{mix}}}$  and let  $\widetilde{E}$  be a set of edges on V, disjoint from E such that each vertex from V belongs to exactly one edge in  $\widetilde{E}$ . Let Y be a weighted random walk on  $(V, E \cup \widetilde{E})$  such that edges in E have weight 1 and edges in  $\widetilde{E}$  have weight  $\theta$ , or more precisely Y is a Markov chain with transition probabilities  $\widetilde{P}(x, y) = \mathbbm{1}_{\{\{x, y\} \in E\}} \frac{1}{\theta + d} + \mathbbm{1}_{\{\{x, y\} \in \widetilde{E}\}} \frac{\theta}{\theta + d}$  for all  $x, y \in V$ .

(i) Are X and Y reversible chains? Find their invariant distributions.

(ii) Prove that if X does not exhibit cutoff, then Y also does not exhibit cutoff.

[You can use without proof that for any constant c and any  $f(n) \approx \frac{1}{g(n)}$  such that  $f(n) \to 0$  as  $n \to \infty$  we have  $(1 - cf(n))^{g(n)} \approx 1$ .]

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(a) State Thomson's principle and define an edge-cutset. Prove the Nash-Williams inequality stating that if  $(\Pi_k)_{k=1}^m$  is a sequence of m disjoint edge-cutsets which separate a and z then

$$R_{\text{eff}}(a,z) \ge \sum_{k=1}^{m} \left( \sum_{e \in \Pi_k} c(e) \right)^{-1}$$

State the commute time identity.

(b) Let  $B = ([0, n] \times [0, n^2]) \cap \mathbb{Z}^2$  be the set of points with integer coordinates in a box of side lengths n and  $n^2$  and let G be the graph with vertex set B and edges between  $(x_1, x_2)$  and  $(y_1, y_2)$  if and only if  $|x_1 - y_1| + |x_2 - y_2| = 1$ . Let X be a simple random walk on G starting from (0, 0) and let  $T_{(n, n^2)}$  be the first time X visits the vertex  $(n, n^2)$ . Find the order of  $\mathbb{E}_{(0,0)}[T_{(n, n^2)}]$ .

### END OF PAPER