

MAT3

MATHEMATICAL TRIPOS **Part III**

Wednesday, 7 June, 2023 9:00 am to 11:00 am

PAPER 224

INFORMATION THEORY

Before you begin please read these instructions carefully

Candidates have **TWO HOURS** to complete the written examination.

Attempt no more than **THREE** questions.

There are **FOUR** questions in total.

The questions carry equal weight.

STATIONERY REQUIREMENTS

Cover sheet

Treasury tag

Script paper

Rough paper

SPECIAL REQUIREMENTS

None

**You may not start to read the questions
printed on the subsequent pages until
instructed to do so by the Invigilator.**

1

- (a) Suppose X, Y are independent RVs with values in $A = \{0, 1, \dots\}$. Show that $H(X + Y) \geq \max\{H(X), H(Y)\}$.
- (b) Suppose X has PMF P_X and mean $\mu > 0$ on A and let Z be an independent geometric with mean μ , i.e., with parameter $p = \frac{1}{1+\mu}$ and PMF $P_Z(k) = p(1-p)^k$, $k \geq 0$. Using the result of part (a) show that,

$$D(P_X \| P_Z) \leq 2d_R(X, Z),$$

where $d_R(X, Z)$ is the *Ruzsa distance* between X and Z :

$$d_R(X, Z) := H(X - Z) - \frac{1}{2}H(X) - \frac{1}{2}H(Z).$$

- (c) Suppose X has PMF P_X on A with $P_X(0) = 0$, let $q := \sum_{k=0}^{\infty} \min\{P_X(k), P_X(k+1)\}$, and write P_{X-1} for the PMF of $X - 1$ on A . Show that,

$$\|P_X - P_{X-1}\|_{\text{TV}} = 2(1 - q).$$

- (d) With X as in part (c), use the result of part (c) to show that:

$$2^{-2H(X)+1} \leq (\log_e 2)D(P_X \| P_{X-1}).$$

Hint. Find an upper bound for q and a lower bound for $H(X)$, both in terms of $\max_k P_X(k)$.

2 Let A be a finite alphabet and Q be a PMF on A .

- (a) Let x_1^n be a string in A^n . Define its type $P = \hat{P}_{x_1^n}$, define the type class $T(P)$, and state an upper bound for the probability $Q^n(T(P))$.

Let B be an arbitrary nonempty subset of A^n . In the next five parts you will establish a corresponding upper bound for $Q^n(B)$.

- (b) Suppose that $Y_1^n = (Y_1, Y_2, \dots, Y_n)$ are independent and identically distributed with each $Y_i \sim Q$, and that X_1^n are distributed as Y_1^n conditional on $Y_1^n \in B$. Write down a formula for the joint PMF P_n of X_1^n .
- (c) Let J be uniformly distributed on $\{1, 2, \dots, n\}$, independent of Y_1^n . Derive a formula for the PMF \bar{P} of X_J in terms of the types $\hat{P}_{x_1^n}$ of the strings $x_1^n \in B$.
- (d) Show that $H(X_1^n) \leq nH(\bar{P})$.

- (e) Show that:

$$\sum_{x_1^n \in B} \frac{Q^n(x_1^n)}{Q^n(B)} \log Q^n(x_1^n) = n \sum_{x \in A} \bar{P}(x) \log Q(x).$$

- (f) Using parts (b), (d) and (e) show that: $Q^n(B) \leq 2^{-nD(\bar{P}\|Q)}$.

3

- (a) State Kraft's inequality for prefix-free free codes (C_n, L_n) on A^n for a finite alphabet A .
- (b) State and prove both the direct and converse parts of the codes-distributions correspondence.
- (c) Let $X_1^n = (X_1, \dots, X_n)$ be random variables with values in the finite alphabet A , and let $W(x_1^n)$ denote the "weight" of a string $x_1^n \in A^n$ for some fixed weight function $W : A^n \rightarrow (0, \infty)$. Find the smallest achievable value of the average weighted description length, $\mathbb{E}[W(X_1^n)L_n(X_1^n)]$, among all prefix-free codes, ignoring integer codelength constraints. Describe the length function L_n^* that achieves that minimum.

4

- (a) Suppose X_1, X_2, \dots, X_n are (not necessarily independent) Bernoulli random variables (RVs), and let $S_n = X_1 + \dots + X_n$. State and prove a Poisson approximation bound for the PMF P_{S_n} of S_n . You may assume, without proof, that $D_e(\text{Bern}(q) \parallel \text{Po}(q)) \leq q^2$.
- (b) Let $\{X_i^{(n)}\} = \{(X_1^{(n)}, X_2^{(n)}, \dots, X_n^{(n)}) ; n \geq 1\}$ be a triangular array of independent Bernoulli RVs, where, for each row $n \geq 1$, the RVs $(X_1^{(n)}, X_2^{(n)}, \dots, X_n^{(n)})$ are IID $\text{Bern}(\lambda/n)$ for some fixed $\lambda > 0$ independent of n .

Let P_n denote the joint PMF of $(X_1^{(n)}, X_2^{(n)}, \dots, X_n^{(n)})$, $n \geq 1$. Find a sequence of constants $\{c_n\}$ and a RV Z such that the following version of the asymptotic equipartition property holds in this case: As $n \rightarrow \infty$:

$$-\frac{1}{c_n} \log P_n(X_1^{(n)}, X_2^{(n)}, \dots, X_n^{(n)}) \rightarrow Z \quad \text{in distribution.}$$

You may assume the following without proof: If $\{Z_n\}$ and Z are RVs and $\{a_n\}$ and $\{b_n\}$ are sequences of real numbers such that, as $n \rightarrow \infty$, (i) the PMFs P_{Z_n} of Z_n converge to the PMF P_Z of Z in that $\|P_{Z_n} - P_Z\|_{\text{TV}} \rightarrow 0$, (ii) $a_n \rightarrow 1$, and (iii) $b_n \rightarrow 0$, then $a_n Z_n + b_n \rightarrow Z$ in distribution, as $n \rightarrow \infty$.

END OF PAPER