

MAT3, MAMA

MATHEMATICAL TRIPOS **Part III**

Thursday, 6 June, 2019 1:30 pm to 3:30 pm

PAPER 215

MIXING TIMES OF MARKOV CHAINS

*Attempt no more than **THREE** questions.*

*There are **FOUR** questions in total.*

The questions carry equal weight.

STATIONERY REQUIREMENTS

Cover sheet

Treasury Tag

Script paper

Rough paper

SPECIAL REQUIREMENTS

None

<p>You may not start to read the questions printed on the subsequent pages until instructed to do so by the Invigilator.</p>

1

Consider a deck of n cards labeled 1 to n . The *random to top* card shuffling scheme is the Markov chain, whose state space is the $n!$ possible arrangements of the cards, that at each step picks a card uniformly at random (possibly the top card) and moves it to the top of the deck. The *top to random* card shuffling takes at each step the top card and insert it at a position chosen uniformly at random (possibly the top position).

Show that as a sequence in n both families of chains exhibit cutoff.

(You may use the solution to the coupon-collector problem as given, as well as the fact that shortly before the coupon collector time with a large probability many coupons are still not picked.)

2

(a) Let $d, n \in \mathbb{N}$. Let e_1, \dots, e_d be the standard basis of $(\mathbb{Z}/n\mathbb{Z})^d$. Let $p_1, \dots, p_d \in [0, 1]$. Consider a random walk $X(k)$ on $(\mathbb{Z}/n\mathbb{Z})^d$, which evolves as follows: $X(k+1) - X(k)$ is distributed as $\sum_{i=1}^d e_i \xi_i$, where ξ_1, \dots, ξ_d are independent and ξ_i is equal to either $-1, 0$ or 1 with probability $\frac{1-p_i}{2}, 1/2$ and $p_i/2$, respectively, for each i . Show that the mixing time is $O(n^2 \log(d+1))$, with the implicit constant being independent of the values of p_1, \dots, p_d . Show for $d = 1$ a matching lower bound.

You may use without a proof the fact that for a *simple* random walk on the cycle, if ρ is the hitting time of 0 then $\max_{i \in \mathbb{Z}/n\mathbb{Z}} \mathbb{E}_i[\rho] \leq n^2$.

(b) Let $d, n \in \mathbb{N}$. Consider a random walk on $\{1, \dots, n\}^d$, in which the d coordinates evolve independently, each according to the following transition matrix: for all $\ell \in \{2, \dots, n-1\}$ let $P(\ell, \ell) = 1/2$ and $P(\ell, \ell+1) = 1/3 = 2P(\ell, \ell-1)$, whereas $P(0, 0) = 2/3 = 2P(0, 1)$ and $P(n, n) = 5/6 = 5P(n, n-1)$ (in other words, each coordinate attempts at each step to move up with probability $1/3$ and down with probability $1/6$, but moves away from the interval $\{1, \dots, n\}$ are censored). We allow d to vary with n but require that $\lim_{n \rightarrow \infty} \frac{\log d}{n} = 0$. Show that the mixing time is $O(n)$.

(c) Show that the family of Markov chains from part (b) (indexed by n) exhibits cutoff. You may assume $d = 1$.

3

Let $G = (V, E)$ be a connected n -vertex simple graph (i.e. it has no self-loops nor multiple edges between a pair of vertices). Let $p \in (0, 1]$. Consider the following Markov chain on $\{\pm 1\}^V$: given that the state at time t is $\sigma_t \in \{\pm 1\}^V$ the next state σ_{t+1} is obtained by picking a vertex v uniformly at random, setting $\sigma_{t+1}(u) = \sigma_t(u)$ for all $u \neq v$ and then:

- With probability p set the value of $\sigma_{t+1}(v)$ to be 1 or -1 with equal probability.
- With probability $1 - p$ a neighbor u of v (i.e. $uv \in E$) is picked uniformly at random and then we set $\sigma_{t+1}(v) = \sigma_t(u)$.

Let P be the corresponding transition matrix.

(i) Show that if G is regular (i.e. all vertices have the same degree), then for all $t \geq 0$ we have that

$$\max_{\sigma, \sigma' \in \{\pm 1\}^V} \|P^t(\sigma, \cdot) - P^t(\sigma', \cdot)\|_{\text{TV}} \leq n \left(1 - \frac{p}{n}\right)^t.$$

(ii) Show that even if G is not regular, for all $t \geq 0$ we have that

$$\max_{\sigma, \sigma' \in \{\pm 1\}^V} \|P^t(\sigma, \cdot) - P^t(\sigma', \cdot)\|_{\text{TV}} \leq 2|E| \left(1 - \frac{p}{n}\right)^t.$$

[Hint: the forms of the bounds suggest what method should be used.]

4

(a) State and prove a theorem demonstrating the use of the canonical paths method to obtain a Poincaré inequality.

(b) Consider a Markov chain $(X_t)_{t \geq 0}$ with a stationary distribution π . Let A be a subset of the state space Ω with $\pi(A) \geq 1/2$. Denote its hitting time by $T_A := \inf\{t : X_t \in A\}$. Show that the $1/4$ total-variation mixing time is at least $c \max_{x \in \Omega} \mathbb{E}_x[T_A]$, where c is some positive absolute constant independent of the Markov chain.

END OF PAPER