

MATHEMATICAL TRIPOS **Part III**

Tuesday, 31 May, 2016 9:00 am to 12:00 pm

PAPER 119

INTRODUCTION TO CATEGORY THEORY

*Attempt no more than **FOUR** questions.*

*There are **SIX** questions in total.*

The questions carry equal weight.

STATIONERY REQUIREMENTS

Cover sheet

Treasury Tag

Script paper

SPECIAL REQUIREMENTS

None

<p>You may not start to read the questions printed on the subsequent pages until instructed to do so by the Invigilator.</p>

1 Define the terms *monomorphism*, *strong* (or *extremal*) *monomorphism* and *regular monomorphism*. Show that any regular monomorphism is strong, and that if two subobjects $A' \twoheadrightarrow A$, $A'' \twoheadrightarrow A$ are strong then so is their intersection (= pullback) $A' \cap A'' \twoheadrightarrow A$, if it exists.

We call a morphism *anodyne* if it is both monic and epic, and we call an object B *saturated* if it is injective with respect to the class of anodyne morphisms, i.e. every diagram

$$\begin{array}{ccc} A' & \xrightarrow{\quad} & A \\ & \downarrow & \\ & B & \end{array}$$

can be completed to a commutative triangle. Show that a strong subobject of a saturated object is saturated.

Now suppose that \mathcal{C} is complete and well-powered, and that every object A of \mathcal{C} admits a monomorphism $A \twoheadrightarrow B$ with B saturated. Show that every object admits an anodyne morphism to a saturated object [*hint: consider the smallest strong subobject of B which contains A*]. Deduce that the full subcategory \mathcal{S} of saturated objects is reflective in \mathcal{C} . Show also that \mathcal{S} is balanced [*hint: first show that epimorphisms in \mathcal{S} are also epic in \mathcal{C}*].

2 State the Yoneda Lemma. If \mathcal{C} is a small category, show that any functor $F : \mathcal{C} \rightarrow \mathbf{Set}$ may be expressed as a colimit in $[\mathcal{C}, \mathbf{Set}]$ of a diagram of shape $(1 \downarrow F)^{\text{op}}$ whose vertices are representable functors, where 1 denotes a singleton set. [*Hint: observe that a natural transformation $\alpha : F \rightarrow G$ is determined by the family of elements $(\alpha_A(x) \mid A \in \text{ob } \mathcal{C}, x \in FA)$* .]

Now suppose \mathcal{C} has finite limits. Show that the following conditions on F are equivalent:

- (i) F preserves finite limits.
- (ii) For any set S , $(S \downarrow F)$ has finite limits.
- (iii) $(1 \downarrow F)^{\text{op}}$ is filtered.
- (iv) F is expressible as a filtered colimit of representable functors.

[You may assume the result that filtered colimits commute with finite limits in \mathbf{Set} .]

3 Let \mathcal{C} be a category with finite products. An object A of \mathcal{C} is said to be *exponentiable* if the functor $(-)\times A: \mathcal{C} \rightarrow \mathcal{C}$ has a right adjoint $(-)^A$. Show that the class of exponentiable objects of \mathcal{C} is closed under finite products.

Now suppose that \mathcal{C} is complete and locally small, that it has a coseparator S , and that every monomorphism in \mathcal{C} is regular. Show that for every object B of \mathcal{C} there is an equalizer diagram of the form

$$B \rightrightarrows \prod_{i \in I} S \rightrightarrows \prod_{j \in J} S$$

for suitable index sets I and J . Deduce that an object A is exponentiable if and only if the functor $\mathcal{C}(-\times A, S): \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}$ is representable. [*Hint: if the functor $(-)^A$ exists, it preserves limits.*]

4 Explain what is meant by the *monadic length* of an adjunction

$$(F: \mathcal{C} \rightarrow \mathcal{D} \dashv G: \mathcal{D} \rightarrow \mathcal{C}),$$

where \mathcal{D} has reflexive coequalizers. [You should write down the definitions of the Eilenberg–Moore comparison functor and its left adjoint, but you need not verify their properties.]

For each natural number n , let \mathcal{C}_n denote the category whose objects are sets A equipped with n partial unary operations $\alpha_1, \alpha_2, \dots, \alpha_n$, such that $\alpha_1(a)$ is defined for all $a \in A$, and for $i > 1$ $\alpha_i(a)$ is defined if and only if $(\alpha_{i-1}(a)$ is defined and) $\alpha_{i-1}(a) = a$, and whose morphisms $(A, \alpha_1, \dots, \alpha_n) \rightarrow (B, \beta_1, \dots, \beta_n)$ are functions $f: A \rightarrow B$ such that $f(\alpha_i(a)) = \beta_i(f(a))$ whenever $\alpha_i(a)$ is defined, for each i . Show that the forgetful functor $G_n: \mathcal{C}_{n+1} \rightarrow \mathcal{C}_n$ has a left adjoint F_n [*hint: $F_n(A, \alpha_1, \dots, \alpha_n)$ may be taken to have underlying set $(A \times \{0\}) \cup (A_n \times \mathbb{N})$, where $A_n = \{a \in A \mid \alpha_n(a) = a\}$.]*]

Show also that, for each $m > n$, the composite adjunction

$$\mathcal{C}_n \begin{array}{c} \xrightarrow{F_n} \\ \xleftarrow{G_n} \end{array} \mathcal{C}_{n+1} \begin{array}{c} \xrightarrow{F_{n+1}} \\ \xleftarrow{G_{n+1}} \end{array} \mathcal{C}_{n+2} \cdots \mathcal{C}_{m-1} \begin{array}{c} \xrightarrow{F_{m-1}} \\ \xleftarrow{G_{m-1}} \end{array} \mathcal{C}_m$$

induces the same monad (up to isomorphism) on \mathcal{C}_n as $\mathcal{C}_n \rightleftarrows \mathcal{C}_{n+1}$. Show that G_n creates coequalizers of G_n -split pairs, and deduce that the composite adjunction displayed above has monadic length $m-n$.

5 Explain carefully what is meant by the assertion that limits of shape I commute with colimits of shape J in a category \mathcal{C} .

Let G be a group, considered as a category with one object, and let A be a G -set, regarded as a functor $G \rightarrow \mathbf{Set}$. Show that $\lim_G A$ and $\operatorname{colim}_G A$ may be identified respectively with the set of G -fixed elements of A and the set of G -orbits.

Hence show that

(a) if G and H are finite groups of coprime orders, then limits of shape G commute with colimits of shape H in \mathbf{Set} ;

(b) if G and H have a nontrivial common quotient group K , then limits of shape G do not commute with colimits of shape H in \mathbf{Set} .

[Hint for (b): let G and H act on (the underlying set of) K by $(g, k) \mapsto q(g).k$ and $(h, k) \mapsto k.r(h)^{-1}$, where q and r are the quotient maps.]

6 Define the notion of *abelian category*, and prove from your definition that epimorphisms in an abelian category are stable under pullback. [Standard results on additive categories may be assumed.]

Now suppose given an exact sequence $0 \rightarrow A \xrightarrow{f} B \xrightarrow{g} C \rightarrow 0$ and a pullback square

$$\begin{array}{ccc} B' & \xrightarrow{g'} & C' \\ \downarrow h & & \downarrow k \\ B & \xrightarrow{g} & C \end{array}$$

in an abelian category. Show that there is a morphism $f': A \rightarrow B'$ making the sequence $0 \rightarrow A \xrightarrow{f'} B' \xrightarrow{g'} C' \rightarrow 0$ exact.

Deduce that if $0 \rightarrow K \rightarrow P \rightarrow C \rightarrow 0$ and $0 \rightarrow K' \rightarrow P' \rightarrow C \rightarrow 0$ are exact sequences with P and P' projective, then $K \oplus P' \cong K' \oplus P$. [Hint: show that they are both isomorphic to the pullback of $P \rightarrow C$ and $P' \rightarrow C$.]

END OF PAPER