

MATHEMATICAL TRIPOS Part III

Monday, 9 June, 2014 1:30 pm to 3:30 pm

PAPER 63

ADVANCED QUANTUM INFORMATION THEORY

*Attempt no more than **TWO** questions.*

*There are **FOUR** questions in total.*

The questions carry equal weight.

STATIONERY REQUIREMENTS

Cover sheet

Treasury Tag

Script paper

SPECIAL REQUIREMENTS

None

<p>You may not start to read the questions printed on the subsequent pages until instructed to do so by the Invigilator.</p>

1 A *stoquastic circuit* is a quantum circuit consisting of 2- and 3-qubit gates which are all permutation matrices in the computational basis.

The input to the stoquastic circuit is always a computational basis state, together with a set A_0 of ancilla qubits each in the $|0\rangle$ state, and a set A_+ of ancilla qubits each in the $|+\rangle$ state, where we define the single qubit states $|+\rangle = \frac{1}{\sqrt{2}}(|0\rangle + |1\rangle)$ and $|-\rangle = \frac{1}{\sqrt{2}}(|0\rangle - |1\rangle)$.

The output of the stoquastic circuit is the result of measuring the first qubit in the $\{|+\rangle, |-\rangle\}$ basis.

The complexity class StoqMA is the class of decision problems for which there exists a polynomial-sized *stoquastic* verifier circuit U acting on a computational basis state $|w\rangle$ such that:

$$\begin{cases} \exists |w\rangle : \Pr(U \text{ outputs "+" on input } |w\rangle) \geq \frac{2}{3} & \text{YES instance} \\ \forall |w\rangle : \Pr(U \text{ outputs "+" on input } |w\rangle) \leq \frac{1}{3} & \text{NO instance} \end{cases}$$

A *stoquastic Hamiltonian* is a Hamiltonian whose off-diagonal matrix elements in the computational basis are all ≤ 0 .

- (a) Use the Kitaev construction with non-local clock to write down a Hamiltonian H encoding a stoquastic verifier circuit (*without* any penalty term on the output of the circuit).

Show that H is stoquastic.

- (b) Prove that the 0-energy eigenstates of H are computational history states for the stoquastic verifier circuit.

You may use without proof the fact that the matrix $E = \frac{1}{2} \sum_{t=0}^T (|t\rangle - |t+1\rangle)(\langle t| - \langle t+1|)$ has eigenvalues $\lambda_k = 1 - \cos q_k$ where $q_k = \frac{\pi k}{T+1}$, $k = 0, \dots, T$, and that the eigenvector corresponding to λ_0 is $\frac{1}{\sqrt{T+1}} \sum_{t=0}^T |t\rangle$.

- (c) Prove that the spectral gap $\Delta(H) = \Omega(T^{-3})$.

You may use without proof the following variant of Kitaev's geometrical lemma (where $\text{supp } X$ denotes the support of an operator X):

Let $A \geq 0$, $B \geq 0$ be positive-semidefinite operators satisfying $\lambda_{\min}(A|_{\text{supp } A}) \geq \mu$ and $\lambda_{\min}(B|_{\text{supp } B}) \geq \mu$. The smallest non-zero eigenvalue of $A + B$ satisfies

$$\lambda_{\min}\left((A + B)|_{\text{supp}(A+B)}\right) \geq 2\mu \sin^2 \frac{\theta}{2}$$

where $\theta := \max_{\substack{|\xi\rangle \in \ker A \\ |\eta\rangle \in |B\rangle \\ |\eta\rangle \perp K}} \langle \eta | \xi \rangle$ with $K = \ker A \cap \ker B$.

You may also use without proof all the properties of the matrix E given in part (b), and the relation $\sin^2 \frac{\theta}{2} \geq \frac{1 - \cos^2 \theta}{8 \cos^2 \theta}$ valid for $0 \leq \theta \leq \frac{\pi}{4}$.

- (d) Let $H' = H + \delta H_{\text{out}}$, where $H_{\text{out}} := |-\rangle\langle -|_1 \otimes |T\rangle\langle T|$ and $\delta > 0$ is a constant.

By considering H' , or otherwise, prove that the local Hamiltonian problem for stoquastic Hamiltonians (with non-local clock) is StoqMA-hard.

You may use without proof the following perturbation bound for arbitrary Hamiltonians H and V , valid for any $\delta < \Delta(H)$:

$$\lambda_{\min}(H + \delta V) = \delta \min_{\psi \in \ker H} \langle \psi | H_{\text{out}} | \psi \rangle + O(\delta^2).$$

[*Hint: Show first that for a computational history state $|\psi\rangle$ for circuit input $|w\rangle$, $\langle \psi | H_{\text{out}} | \psi \rangle = \frac{1}{T+1} \left(1 - \Pr(U \text{ outputs “+” on input } |w\rangle) \right)$.]*

2

- State the definition of the complexity class BQP, and explain why the class does not depend on the precise values chosen for the probabilities appearing in the definition.
- Define the complexity class C to be the class of decision problems for which there exists a classical algorithm that, for problem size n , solves the problem using an arbitrary number of basic arithmetic operations on real numbers (addition, subtraction, multiplication), but only ever stores at most $\text{poly}(n)$ real numbers in memory.

Show that $\text{BQP} \subseteq \text{C}$.

[*Hint: The probability of measuring $\Pi_1^{(1)}$ on the output of circuit U running on input $|x\rangle$ can be computed as*

$$\sum_{i_2, i_3, \dots, i_n=0,1} \left| \langle 1 | \langle i_2 | \cdots \langle i_n | U | x \rangle \right|,$$

and the product $U_1 U_2$ can be rewritten as

$$U_1 U_2 = \sum_{i_1, i_2, \dots, i_n=0,1} U_1 |i_1\rangle |i_2\rangle \cdots |i_n\rangle \langle i_1 | \langle i_2 | \cdots \langle i_n | U_2. \quad]$$

- State the definition of the complexity class QMA, and show that $\text{QMA} \subseteq \text{C}$.

You may use without proof the following inequality for the maximum eigenvalue $\lambda_{\max}(X)$ of an n -qubit operator X :

$$\ln \lambda_{\max}(X) \leq \frac{1}{d} \ln \text{Tr}(X^d) \leq \ln \lambda_{\max}(X) + \frac{n}{d} \ln 2.$$

[*Hint: Consider computing the maximum probability (over all witnesses) that the verifier circuit outputs “1”.*]

3 For a system of qudits, let A_X and B_Y be operators that act non-trivially only on qudits in the sets X and Y . $A_X(t)$ denotes the time-evolution of A_X in the Heisenberg picture.

Let Z range over k -element subsets of the qudits, and let $H = \sum_Z h_Z$ be a k -local Hamiltonian where h_Z acts non-trivially only on the subset Z .

- (a) Suppose that for each qudit i in the system, $\sum_{Z:Z\ni i} \|h_Z\| \leq se^{-\mu}$ for some constants $\mu, s > 0$ (where $\|X\|$ denotes the operator norm of X).

Starting from the integral equation

$$C_A(Y, t) \leq C_A(Y, 0) + 2 \sum_{Z:Y \cap Z \neq \emptyset} \|h_Z\| \int_0^t C(Z, s) ds$$

where (for operators O_Y acting non-trivially only on Y)

$$C_A(Y, t) := \sup_{O_Y} \frac{\|[A_X(t), O_Y]\|}{\|O_Y\|},$$

prove the Lieb-Robinson bound:

$$\|[A_X(t), B_Y]\| \leq 2 \|A_X\| \|B_Y\| \min(|X|, |Y|) e^{-\mu d(X, Y)} (e^{2kst} - 1),$$

where $d(X, Y)$ denotes interaction distance.

- (b) In this part of the question, you may use without proof the following corollary of the Lieb-Robinson bound:

$$\|A_X(t) - A_{X'}(t)\| \leq \mu vt |X| \|A_X\| e^{-\mu l/2},$$

where $A_{X'}(t)$ acts non-trivially only on the set of qudits $X' = \{i : d(i, X) \leq vt + l\}$, and $v = 4ks/\mu$ is the Lieb-Robinson velocity.

Let $L = \text{diam}(S)$ be the total diameter of the many-body system on which H acts (where $\text{diam}(X) := \max_{i, j \in X} d(i, j)$).

A state $|\psi_1\rangle$ has *topological quantum order* if there exists another state $|\psi_2\rangle$ with the following properties:

- i. $\langle \psi_2 | \psi_1 \rangle = 0$
- ii. For any A_X with $\|A_X\| \leq 1$ and $\text{diam}(X) \leq cL$:
 $|\langle \psi_1 | A_X | \psi_1 \rangle - \langle \psi_2 | A_X | \psi_2 \rangle| \leq \epsilon$

for some constants $c, \epsilon < 1$.

Let $|\psi_0\rangle$ be a state which evolves under H into $|\psi_0(t)\rangle = |\psi_1\rangle$.

Prove that if $t = cL$ for a sufficiently small constant $c > 0$, then $|\psi_0\rangle$ itself must be topologically ordered. (You are *not* required to find an explicit expression for c .)

[Hint: Consider a second state $|\tilde{\psi}_0\rangle$ that evolves into $|\psi_2\rangle$.]

4 In this question, you may use without proof the fact that there exists a function $w(t)$ with Fourier transforms $\hat{w}(E)$ satisfying:

$$w(t) \geq 0, \quad \hat{w}(E) = 0 \text{ for } E \geq \Delta, \quad \hat{w}(0) = 1, \quad \text{and} \\ \int_T^\infty w(t) dt = O\left((\ln \beta T)^2 e^{-\beta T / (\ln \beta T)^2}\right) \text{ for some constants } c, \beta > 0.$$

For a system of qudits, let Z range over k -element subsets of the qudits and let $H = \sum_Z h_Z$ be a k -local Hamiltonian, where h_Z acts non-trivially only on the subset Z .

You may assume that H satisfies the Lieb-Robinson bound

$$\|[A_X(t), B_Y]\| \leq 2 \|A_X\| \|B_Y\| \min(|X|, |Y|) e^{-\mu d(X,Y)} (e^{2kst} - 1)$$

for operators A_X and B_Y that act non-trivially only on qudits in the subsets X and Y . Here, $A_X(t)$ denotes the time-evolution of A_X under H in the Heisenberg picture, and $d(X, Y)$ denotes the interaction distance between subsets X and Y .

- (a) Assume that H has a unique ground state $|\phi_0\rangle$, eigenstates $|\phi_i\rangle$, and spectral gap $\Delta > 0$.

Explain how the Fourier transform and filtering technique can be used to construct an operator $A^{(Z)}$ with matrix elements:

- i. $\langle \phi_0 | A^{(Z)} | \phi_0 \rangle = \langle \phi_0 | h_Z | \phi_0 \rangle$,
- ii. $\langle \phi_0 | A^{(Z)} | \phi_i \rangle = \langle \phi_i | A^{(Z)} | \phi_0 \rangle = 0$ for $i > 0$.

- (b) Using the Fourier transform and filtering technique, or otherwise, prove that H can be rewritten as $H = \sum_Z A^{(Z)}$ with $[A^{(Z)}, P_0] = 0$, where $P_0 := |\phi_0\rangle\langle\phi_0|$.

(The operators $A^{(Z)}$ may act on all the qudits in the system.)

- (c) Let

$$A^{(Z,d)} := \int_{-\infty}^{\infty} w(t) (e^{itH_d} h_Z e^{-itH_d} - e^{itH_{d-1}} h_Z e^{-itH_{d-1}}) dt,$$

where $H_d := \sum_{Y: d(Z,Y) \leq d} h_Y$, so that $H_0 = h_Z$ and $H_D = H$.

By summing $A^{(Z,d)}$ over d , or otherwise, show that $A^{(Z)} = h_Z + \sum_{d=1}^D A^{(Z,d)}$ for suitably chosen $A^{(Z)}$ from part (b).

- (d) Assume that $\sum_{Y:d(Z,Y)=d} \|h_Y\| \min(|Y|, |Z|) = O(d^\alpha)$ for some constant $\alpha > 0$. Starting from the relation

$$\|e^{itH_d} h_Z e^{-itH_d} - e^{itH_{d-1}} h_Z e^{-itH_{d-1}}\| \leq \sum_{Y:d(Z,Y)=d} \int_0^t \|[h_Z(\tau), h_Y]\| d\tau$$

where $h_Z(\tau) := e^{i\tau H_d} h_Z e^{-i\tau H_d}$, prove that

$$\|e^{itH_d} h_Z e^{-itH_d} - e^{itH_{d-1}} h_Z e^{-itH_{d-1}}\| \leq O\left(\|h_Z\| t d^\alpha e^{-\mu d + 2kst}\right).$$

- (e) Using part (d), or otherwise, prove that $\|A^{(Z,d)}\| \leq O\left(\|h_Z\| d^{\alpha+1} e^{-cd/\ln^2 cd}\right)$ for some constant $c > 0$. (You are *not* required to find an explicit expression for c .)

END OF PAPER